

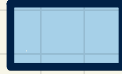
Lecture 8

Part A

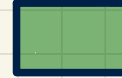
Recursion - Basics: Thinking Recursively, Call Stack

Solving a Problem Recursively

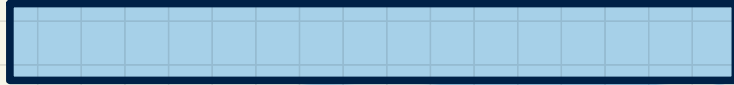
Given a **small** problem:



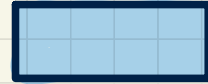
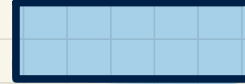
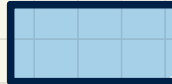
Solve it **directly**:



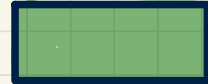
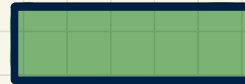
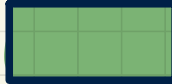
Given a **big** problem:



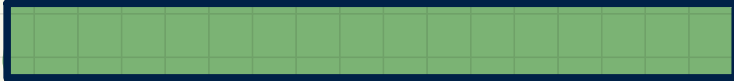
Divide it into **smaller** problems:



Assume solutions to **smaller** problems:



Combine solutions to **smaller** problems:



```
m(i) {  
  if(i == ...) { /* base case: do something directly */ }  
  else {  
    → m(j); /* recursive call with strictly smaller value */  
  }  
} → recursive call of m
```

$m(100)$
↳ $m(99)$
↓
 $j < 100$

Recursive Solution: factorial

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n-1) \cdot (n-2) \cdot \dots \cdot 3 \cdot 2 \cdot 1 & \text{if } \underline{n \geq 1} \end{cases}$$

$$\underline{5!} = 5 * \underline{4 * 3 * 2 * 1}$$

↳ 4!

$$= 5 * 4!$$

$$= 5 * \boxed{(5-1)!}$$

↳ strictly smaller problem size

$$\boxed{n-1 < n}$$

$$\underline{n!} = n * \underline{(n-1) * (n-2) * \dots * 2 * 1}$$

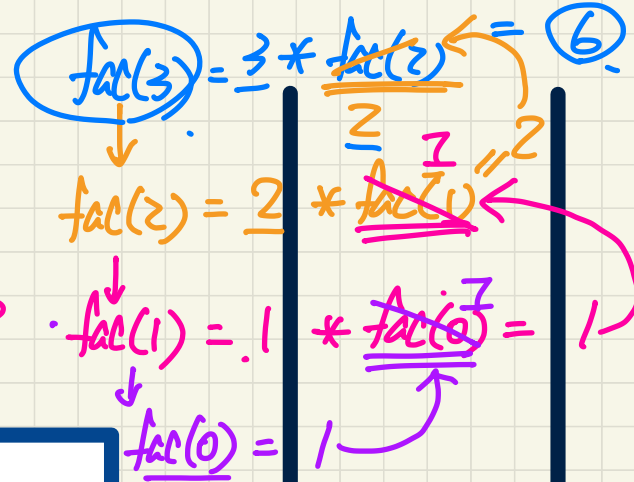
↳ (n-1)!

$$= n * \boxed{(n-1)!}$$

Recursive Solution in Java: factorial

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n-1)! & \text{if } n \geq 1 \end{cases}$$

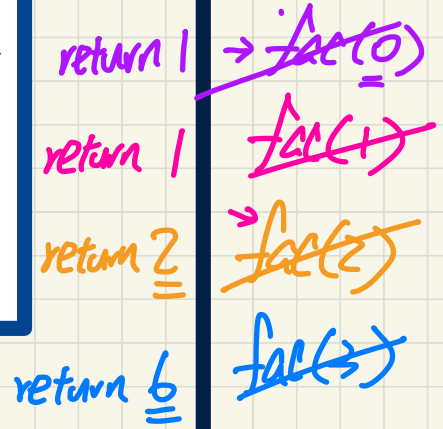
→ strictly smaller problem
→ base case
→ recursive case



```

int factorial (int n) {
    int result;
    if (n == 0) { /* base case */ result = 1; }
    else { /* recursive case */
        result = n * factorial (n - 1);
    }
    return result;
}
    
```

→ strictly smaller problem
→ recursive case
→ 2 * fac(1) = 2 call
→ 1 * fac(0) = 1



Example: factorial(3)

Runtime Stack

Common Errors of Recursion (1)

```
int factorial (int n) {  
    return n * factorial (n - 1);  
}
```

↳ missing base case(s)

Infinite Recursion

fac(3)
↳ fac(2)
↳ fac(1)
↳ fac(0)
↳ fac(-1)
⋮

Common Errors of Recursion (2)

```
int factorial(int n) {  
    if(n == 0) { /* base case */ return 1; }  
    else { /* recursive case */ return n * factorial(n); }  
}
```

Infinite Recursion

fac(3)

↳ fac(3)

↳ fac(3)

↳ fac(3)

⋮

(never able to reach the base case)

problem size for
recursive call is not
strictly smaller.

Recursive Solution: Fibonacci Numbers

$$F = \overset{\cdot}{\circledast}, \overset{\cdot}{\circledast}, \overset{\cdot}{\circledast}, \overset{\cdot}{\circledast}, \overset{\cdot}{\circledast}, \overset{\cdot}{\circledast}, \overset{F_7}{\circledast}, \overset{F_8}{\circledast}, \overset{F_9}{\circledast}, 55, 89, \dots$$

(Note: In the original image, the first two 1s are circled in blue, 13 and 21 are circled in pink, and 34 is circled in green.)

Base Cases

$$F_1 = 1$$

$$F_2 = 1$$

Recursive Cases

$$F_n = F_{n-1} + F_{n-2}$$

(Note: In the original image, F_n is circled in blue, and F_{n-1} and F_{n-2} are circled in green.)

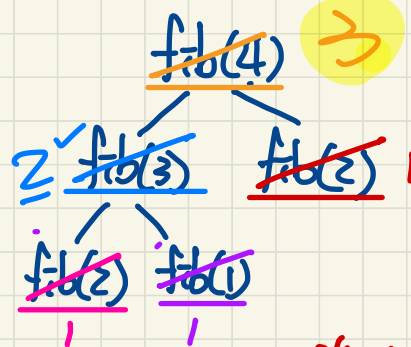
$n > 2$
strictly smaller than n

solved recursively by two recursive calls

$$F_9 = F_7 + F_8$$

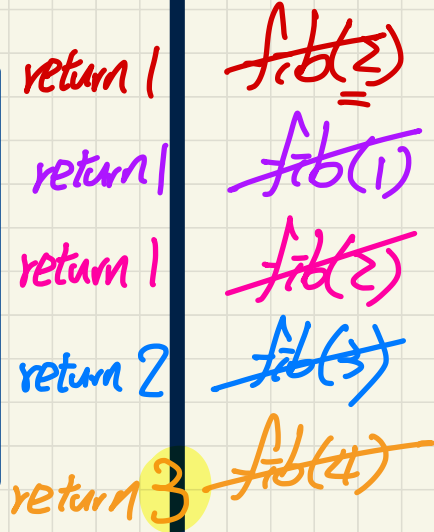
Recursive Solution in Java: Fibonacci Numbers

$$F_n = \begin{cases} 1 & \text{if } n = 1 \\ 1 & \text{if } n = 2 \\ F_{n-1} + F_{n-2} & \text{if } n > 2 \end{cases}$$



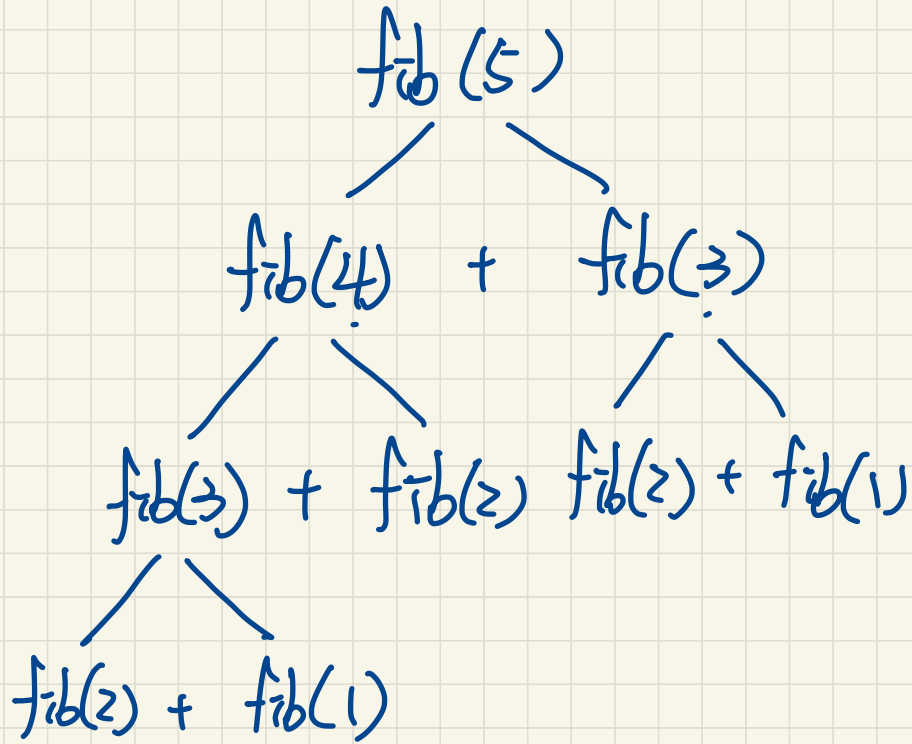
```
int fib (int n) {
    int result;
    if(n == 1) { /* base case */ result = 1; }
    else if(n == 2) { /* base case */ result = 1; }
    else { /* recursive case */
        result = fib (n - 1) + fib (n - 2);
    }
    return result;
}
```

$2 \text{ fib}(3) + \text{fib}(2) = 3$
 $\text{fib}(2) + \text{fib}(1) = 2$



Example: fib(4)

Runtime Stack



Lecture 7

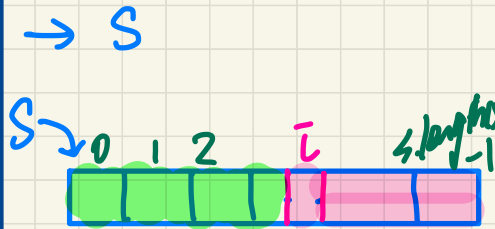
Part B

Recursion - Examples: Recursions on Strings

Use of String

substring(n, m) $\begin{matrix} [n, m) \\ \downarrow \\ [n, m-1] \end{matrix}$ s \rightarrow "abcd"

Assume
i value index
s.substring(0, i) +
s.substring(i, s.length());



```
public class StringTester {  
    public static void main(String[] args) {  
        String s = "abcd";  
        System.out.println(s.isEmpty()); /* false */  
        /* Characters in index range [0, 0) */  
        String t0 = s.substring(0, 0);  
        System.out.println(t0); /* "" */  
        /* Characters in index range [0, 4) */  
        String t1 = s.substring(0, 4);  
        System.out.println(t1); /* "abcd" */  
        /* Characters in index range [1, 3) */  
        String t2 = s.substring(1, 3);  
        System.out.println(t2); /* "bc" */  
        String t3 = s.substring(0, 2) + s.substring(2, 4);  
        System.out.println(s.equals(t3)); /* true */  
        for(int i = 0; i < s.length(); i++) {  
            System.out.print(s.charAt(i));  
        }  
        System.out.println();  
    }  
}
```



Inclusive
Exclusive

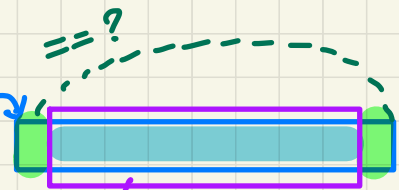
length

s.charAt(0) \rightarrow 'a'
s.charAt(s.length()-1) \rightarrow 'd'

Recursions on Strings

Palindrome

→ "racecar" → T
"aracecars" → F
"raceacar" → F



strictly smaller problem

Reversal

"abcd"

"dcba"

reverse of

strictly smaller problem

solution strictly to smaller prob.

Number of Occurrences

"abca"

'a'

$$2 = 1 + 1$$

'b'

$$1 = 0 + 1$$

of occurrences of the char in tail of input string.

→ the char equal to the head of string.

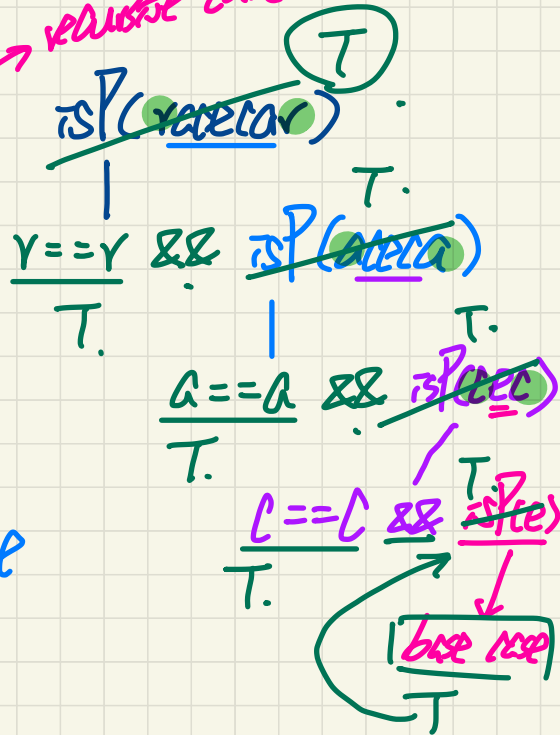
Problem: Palindrome



```

boolean isPalindrome (String word) {
    if (word.length() == 0 || word.length() == 1) {
        /* base case */
        return true;
    }
    else {
        /* recursive case */
        char firstChar = word.charAt(0);
        char lastChar = word.charAt(word.length() - 1);
        String middle = word.substring(1, word.length() - 1);
        return
            firstChar == lastChar
            /* See the API of java.lang.String.substring. */
            && isPalindrome (middle);
    }
}
    
```

recursive case

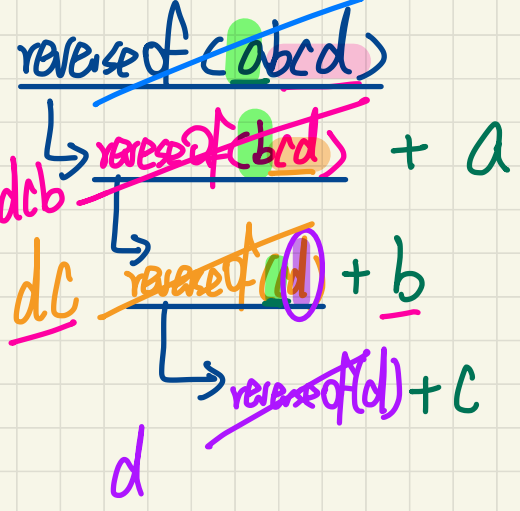


recursive call to solve a subproblem with strictly smaller size

Problem: Reverse of a String

base cases

dcba

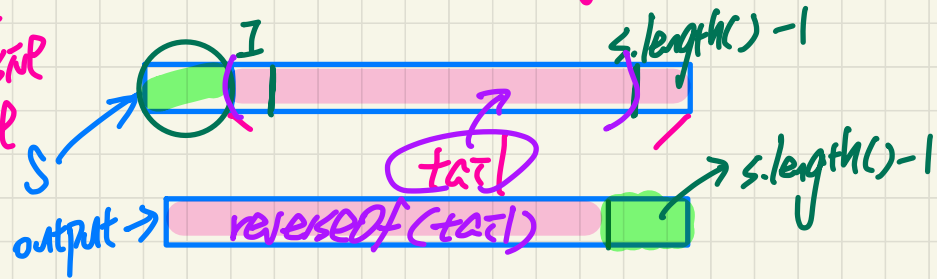


```
String reverseOf (String s) {  
    if(s.isEmpty()) { /* base case 1 */  
        return "";  
    }  
    else if(s.length() == 1) { /* base case 2 */  
        return s;  
    }  
    else { /* recursive case */  
        String tail = s.substring(1, s.length());  
        String reverseOfTail = reverseOf (tail);  
        char head = s.charAt(0);  
        return reverseOfTail + head;  
    }  
}
```

→

↪ recursive call to solve a strictly smaller problem.

RECURSIVE CASE



Problem: Number of Occurrences

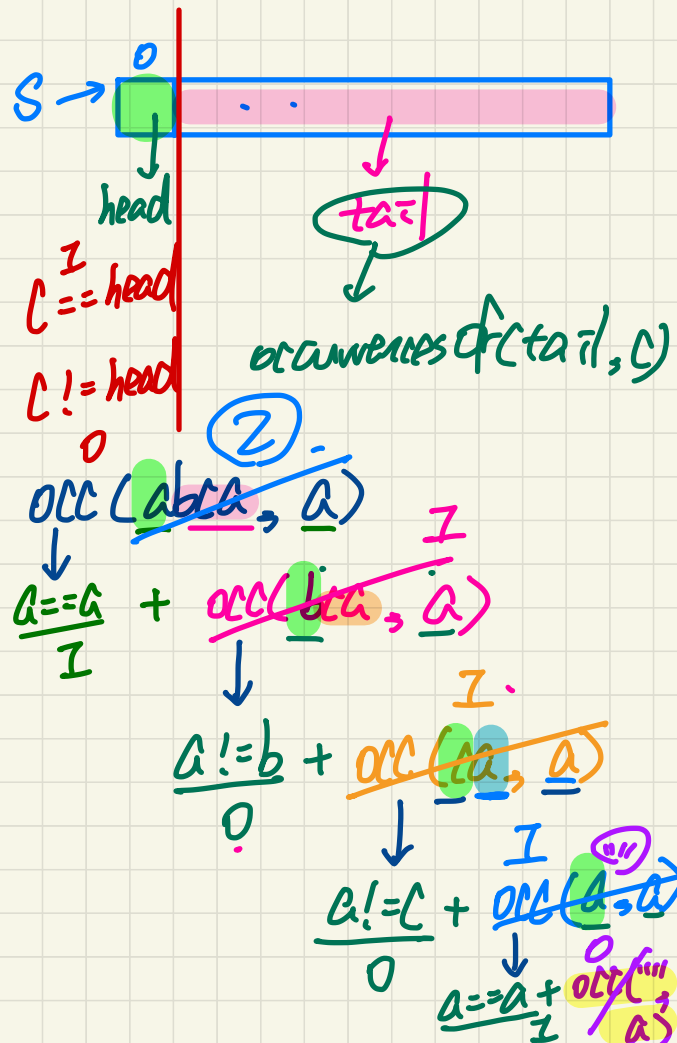
```

int occurrencesOf (String s, char c) {
    if (s.isEmpty()) {
        /* Base Case */
        return 0;
    }
    else {
        /* Recursive Case */
        char head = s.charAt(0);
        String tail = s.substring(1, s.length());
        if (head == c) {
            return 1 + occurrencesOf (tail, c);
        }
        else {
            return 0 + occurrencesOf (tail, c);
        }
    }
}
    
```

→ base case

← recursive case

what if s is "a"?
↳ ""



Lecture 7

Part C

Recursion - Examples: Recursions on Arrays

Recursion on an Array: Passing new Sub-Arrays

```
void m(int[] a) {  
    if(a.length == 0) { /* base case */  
    else if(a.length == 1) { /* base case */  
    else {  
        int[] sub = new int[a.length - 1];  
        for(int i = 1; i < a.length; i++) { sub[i-1] = a[i-1]; }  
        m(sub) } }  
}
```

base cases (green arrow pointing to the if/else if block)

recursive case (pink arrow pointing to the else block)

$i-1$ (pink arrow pointing to the index in the for loop)

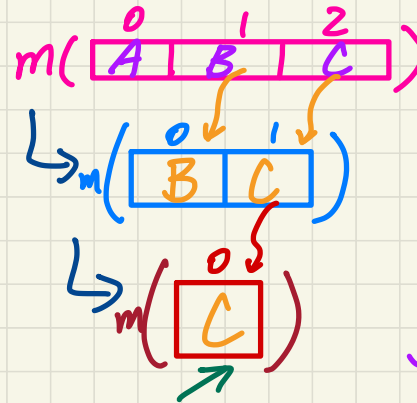
i (pink arrow pointing to the index in the array access)

$sub[0] = a[1]$ (orange arrow pointing to the assignment)

$m(sub)$ (blue arrow pointing to the recursive call)

Say $a_1 = \{\}$ consider $m(a_1)$ → *execute the base case*

Say $a_2 = \{A, B, C\}$, consider $m(a_2)$



not space-efficient (for each v.c., a new array is created)

Recursion on an Array: Passing Same Array Reference

```

void m(int[] a, int from, int to) {
    if (from > to) { /* base case */ }
    else if (from == to) { /* base case */ }
    else { m(a, from + 1, to) } }
    
```

Empty array

array of length 1.

base cases

recursive case

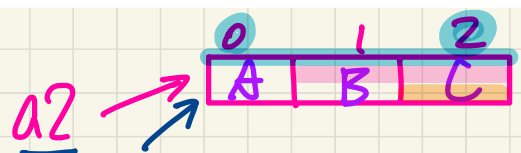
$[0, -1] \rightarrow$ empty range.

Say $a_1 = \{\}$, consider $m(a_1, 0, a_1.length - 1)$

\downarrow min index \downarrow max index
 $\rightarrow m(a_1, 0, -1)$

from to

Say $a_2 = \{A, B, C\}$, consider $m(a_2, 0, a_2.length - 1)$



$m(a_2, 0, 2)$

$\rightarrow m(a_2, 1, 2)$

$\rightarrow m(2, 2)$

strictly smaller problem (last elem in array)

strictly smaller problem (elements from indices 1 to 2)

Problem: Are All Numbers Positive?

```
boolean allPositive(int[] a) {  
    return allPositiveHelper(a, 0, a.length - 1);  
}  
  
boolean allPositiveHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return a[from] > 0;  
    }  
    else { /* recursive case */  
        return a[from] > 0 && allPositiveHelper(a, from + 1, to);  
    }  
}
```

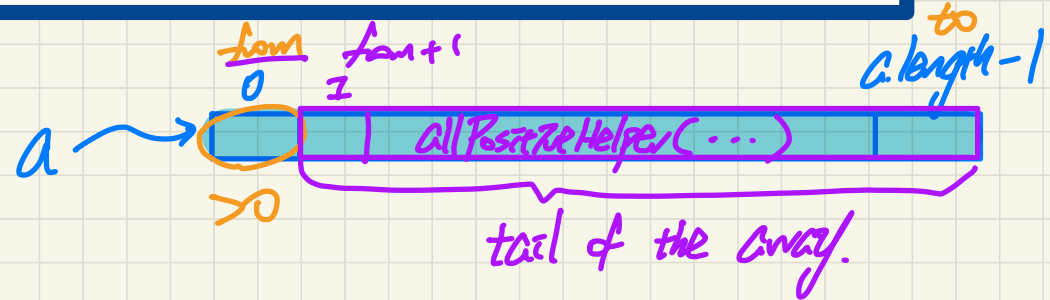
max index
max index
recursive helper method

base cases

empty array

array of length 1

recursive case



Tracing Recursion: allPositive

Say a = {}

allPositive(a)

allPH(a, 0, -1)

```
boolean allPositive(int[] a) {  
    return allPositiveHelper(a, 0, a.length - 1);  
}  
  
boolean allPositiveHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return a[from] > 0;  
    }  
    else { /* recursive case */  
        return a[from] > 0 && allPositiveHelper(a, from + 1, to);  
    }  
}
```

Tracing Recursion: allPositive

Say a = {4}

allPositive(a) ^{4}

allPH(a, 0, 0) ^{a.length - 1}

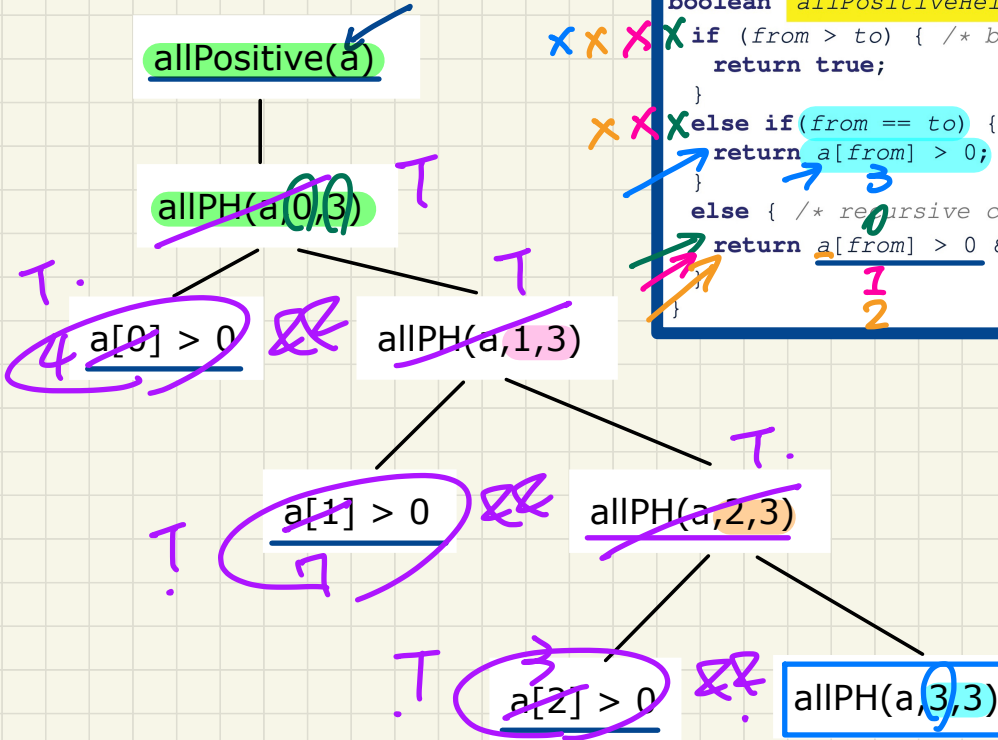
a[0] > 0 ^{True}

4 ⁰ _{from to?}

```
boolean allPositive(int[] a) {  
    return allPositiveHelper(a, 0, a.length - 1);  
}  
  
boolean allPositiveHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return a[from] > 0;  
    }  
    else { /* recursive case */  
        return a[from] > 0 && allPositiveHelper(a, from + 1, to);  
    }  
}
```

Tracing Recursion: allPositive

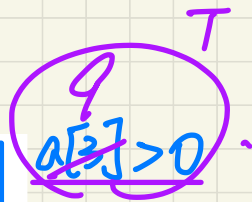
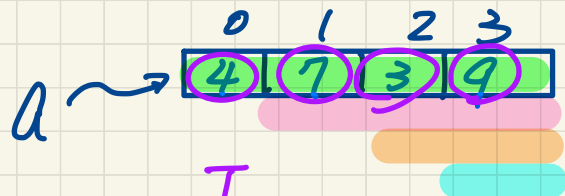
Say a = {4,7,3,9}



```

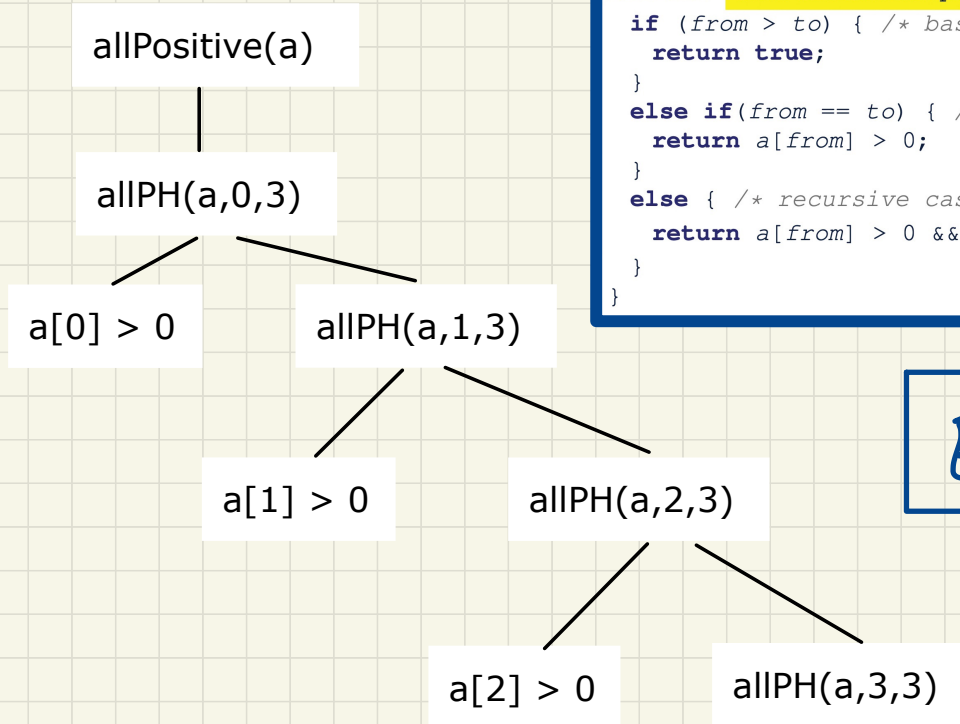
boolean allPositive(int[] a) {
    return allPositiveHelper(a, 0, a.length - 1);
}

boolean allPositiveHelper(int[] a, int from, int to) {
    if (from > to) { /* base case 1: empty range */
        return true;
    }
    else if (from == to) { /* base case 2: range of one element */
        return a[from] > 0;
    }
    else { /* recursive case */
        return a[from] > 0 && allPositiveHelper(a, from + 1, to);
    }
}
  
```



Tracing Recursion: allPositive

Say $a = \{5, 3, -2, 9\}$



```
boolean allPositive(int[] a) {  
    return allPositiveHelper(a, 0, a.length - 1);  
}  
  
boolean allPositiveHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return a[from] > 0;  
    }  
    else { /* recursive case */  
        return a[from] > 0 && allPositiveHelper(a, from + 1, to);  
    }  
}
```

Exercise: Trace!

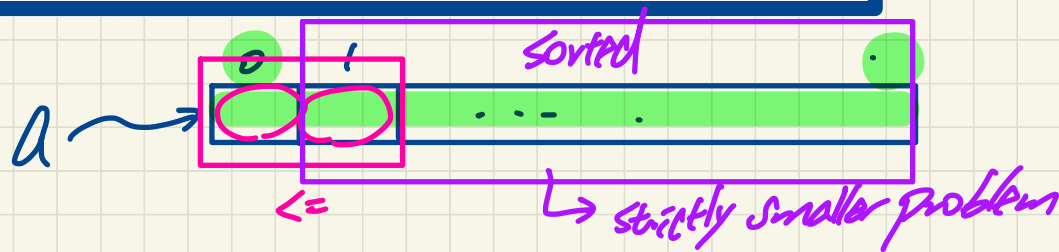
Problem: Are Numbers Sorted?

```
boolean isSorted(int[] a) {  
    return isSortedHelper(a, 0, a.length - 1);  
}  
  
boolean isSortedHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return true;  
    }  
    else {  
        return a[from] <= a[from + 1]  
            && isSortedHelper(a, from + 1, to);  
    }  
}
```

recursive helper method.

base cases

recursive case



Tracing Recursion: isSorted

{}
-1

Say a = {}

isSorted(a) {}

isSH(a, 0, -1)

```
boolean isSorted(int[] a) {  
    return isSortedHelper(a, 0, a.length - 1);  
}  
  
boolean isSortedHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return true;  
    }  
    else {  
        return a[from] <= a[from + 1]  
            && isSortedHelper(a, from + 1, to);  
    }  
}
```

Tracing Recursion: isSorted

Say a = {4}

isSorted(a)

isSH(a,0,0)

return true

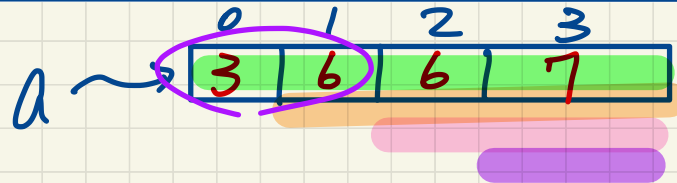
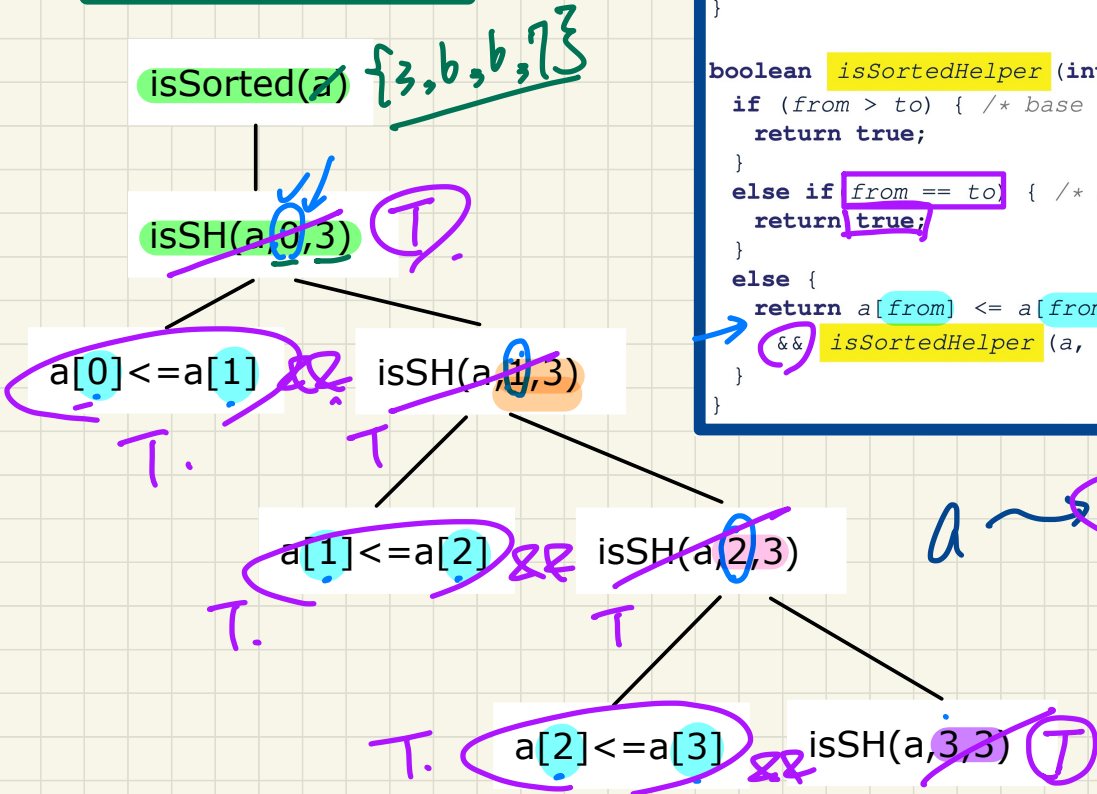
{4}

```
boolean isSorted(int[] a) {  
    return isSortedHelper(a, 0, a.length - 1);  
}  
  
boolean isSortedHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return true;  
    }  
    else {  
        return a[from] <= a[from + 1]  
            && isSortedHelper(a, from + 1, to);  
    }  
}
```

Tracing Recursion: isSorted

Say $a = \{3, 6, 6, 7\}$

```
boolean isSorted(int[] a) {  
    return isSortedHelper(a, 0, a.length - 1);  
}  
  
boolean isSortedHelper(int[] a, int from, int to) {  
    if (from > to) { /* base case 1: empty range */  
        return true;  
    }  
    else if (from == to) { /* base case 2: range of one element */  
        return true;  
    }  
    else {  
        return a[from] <= a[from + 1]  
            && isSortedHelper(a, from + 1, to);  
    }  
}
```



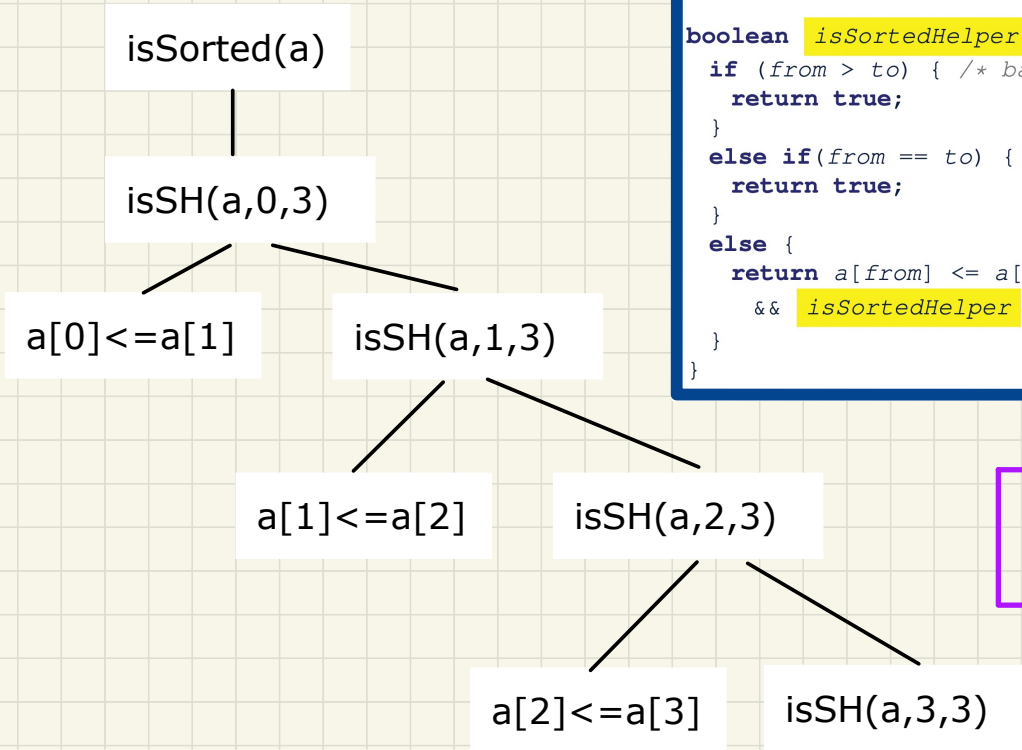
Tracing Recursion: isSorted

Say $a = \{3,6,5,7\}$

(F)

```
boolean isSorted(int[] a) {
    return isSortedHelper(a, 0, a.length - 1);
}

boolean isSortedHelper(int[] a, int from, int to) {
    if (from > to) { /* base case 1: empty range */
        return true;
    }
    else if (from == to) { /* base case 2: range of one element */
        return true;
    }
    else {
        return a[from] <= a[from + 1]
            && isSortedHelper(a, from + 1, to);
    }
}
```



EXERCISE: TRACE